

Optimizing Dynamic Languages for Analytical Workloads

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Analytical Workloads



Analytical

"... concerned with the discovery and communication of meaningful patterns in data"

Analytical Workloads

- > Process large data volumes
- Data often homogeneous, sometimes high-dimensional
- Data mostly read, side-effect free computations

Analytical Workloads



Scenarios

- Decision support
- Statistics
- Model fitting
- Simulations
- **)**

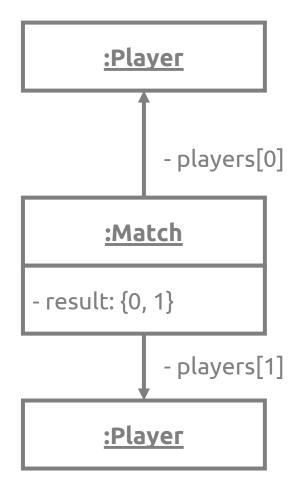
Example

- Scope: ERP web service
- Answer the number of units sold for a specified product

Example: Ranking Players in Games



```
scores = {p: 100 for p in players}
                                 Large data volume
for match in matches_played:
                                 Homogeneous
   pred = predict_result(match, scores)
   delta = 40 * (match.result - pred)
   scores[match.players[0]] += delta
   scores[match.players[1]] -= delta
```



Dynamic Languages



Maintainability Requirements

- Conciseness: code communicates exactly the domain logic
- > Evolution: domain logic should be able to evolve independently from technology and technical code

Dynamic Languages



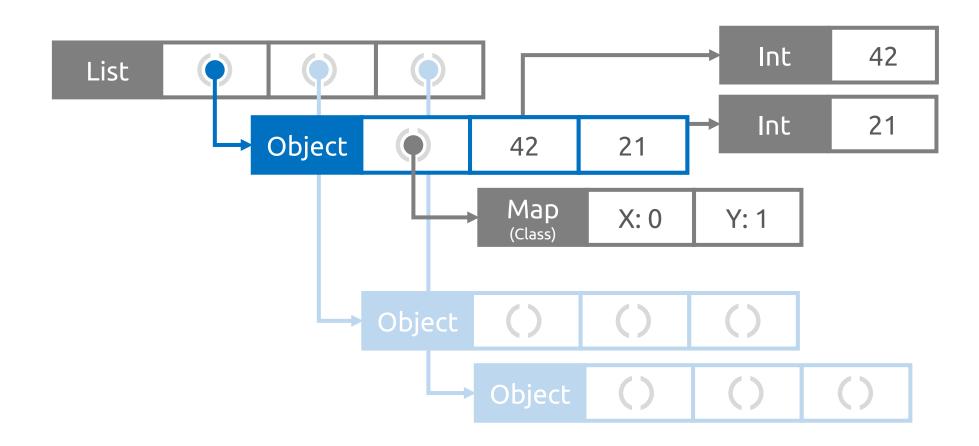
Dynamicity comes at a performance cost

Why?

- Heap structure (objects) build for flexibility
- Analytical data and workloads are quite static

Object Layout





Analytical Databases

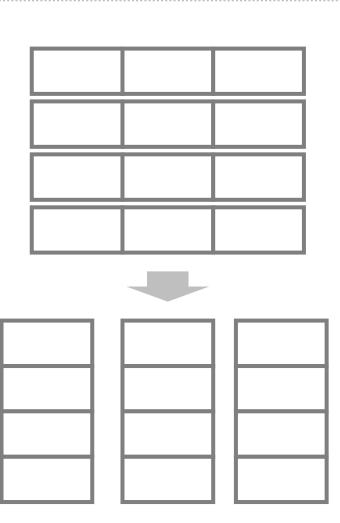


Optimized for

- > Data-intensive queries
- > Reading
- Tough response time requirements

Prevalent Technology

- Main memory based
- Column-oriented storage



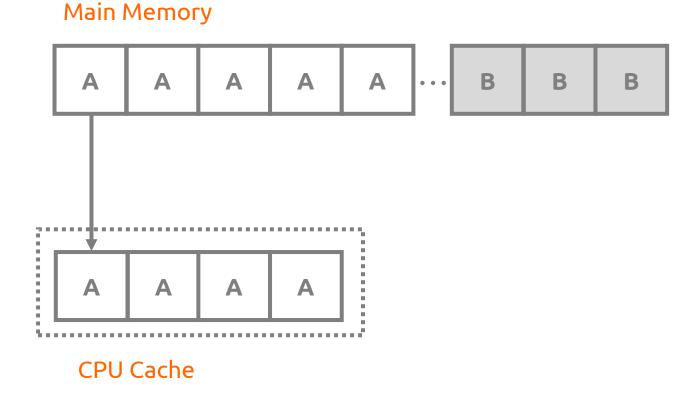
Column-oriented Storage (simplified)



Row-oriented Layout

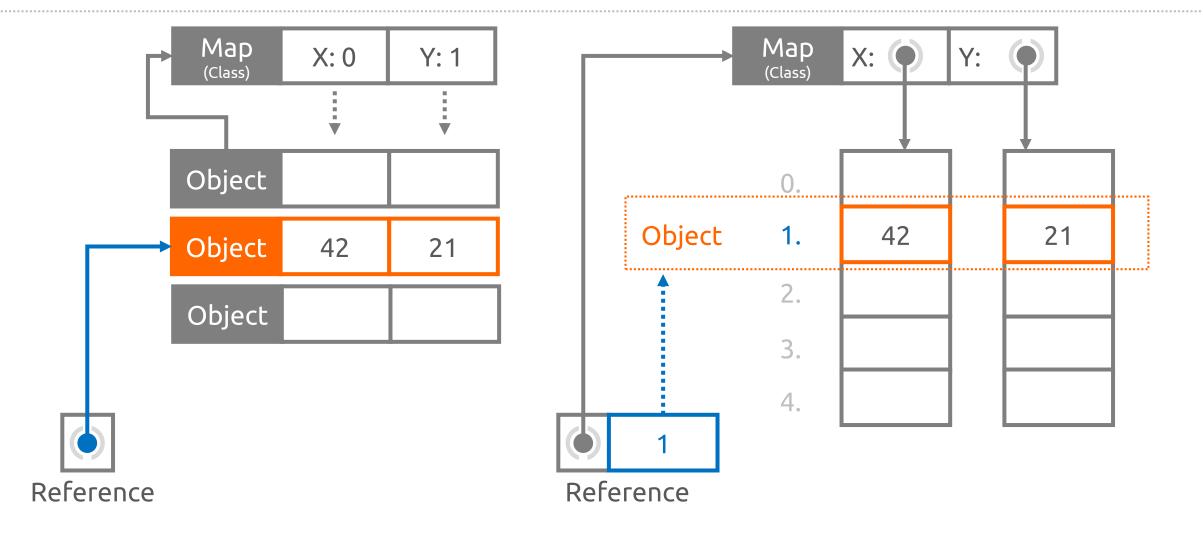
Main Memory "Bottleneck" **CPU Cache**

Column-oriented Layout



Transposing Objects





Object Identity by Proxy

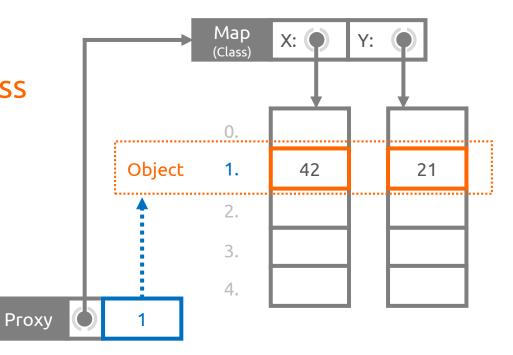


Embedding of (class, offset)-identity into address identity:

Use a proxy object with fields class and offset

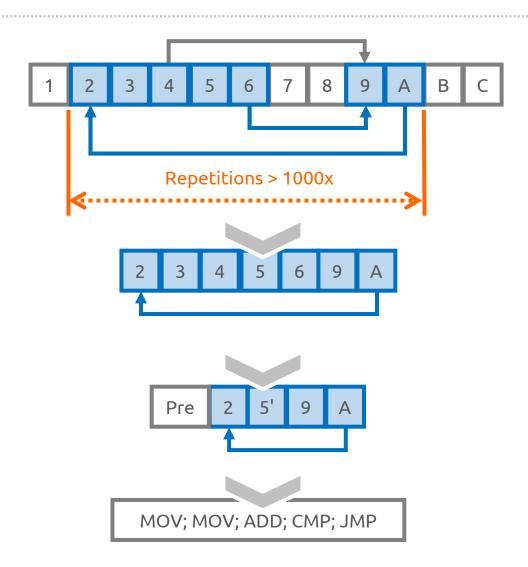
 Dynamically computes memory address on attribute access

» Overhead can be mitigated by Just-in-time (JIT) compilation



Tracing JIT Compilers (Background)





Instruction Stream (Opcodes)

Hot code detected:→ start recording operations

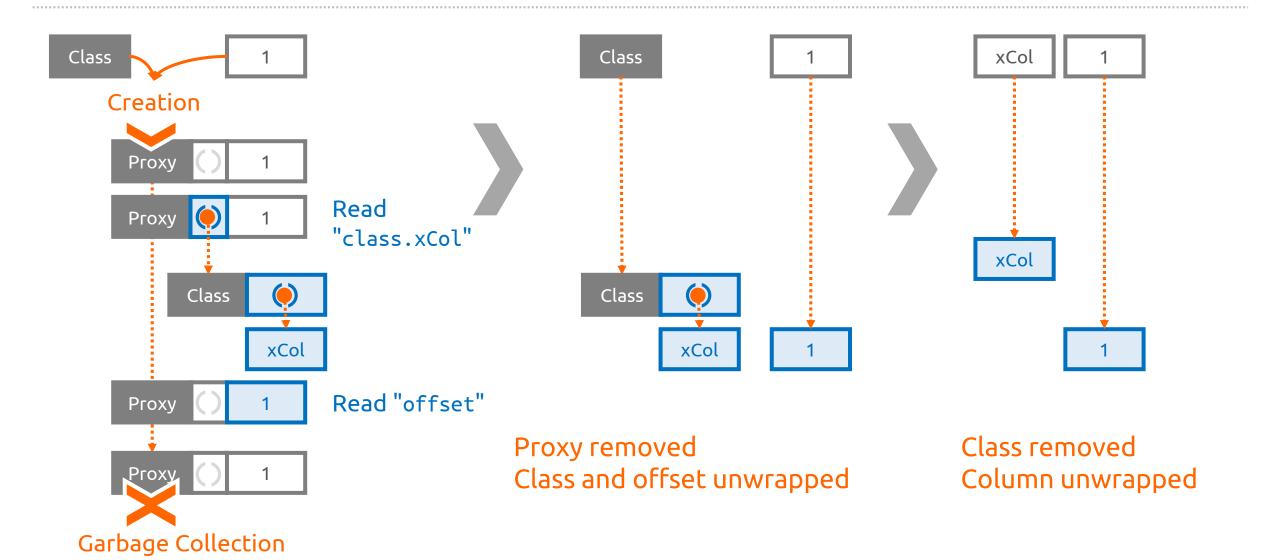
Trace (flat, no branches, only "guards")

Optimized Trace (Loop-invariant code motion, allocation removal, ...)

Native Code

Allocation Removal (Background)





JIT + Columns



» Allocation Removal explodes proxies and classes into offsets and columns

```
for i in LineItem.all():
   total += i.quantity
```

```
n = 0
while n < size:
    i = proxy(LineItem, n)
    col = i.class.columns['quantity']
    offset = i.offset
    value = col[offset]
    total = total + value
    n += 1</pre>
```

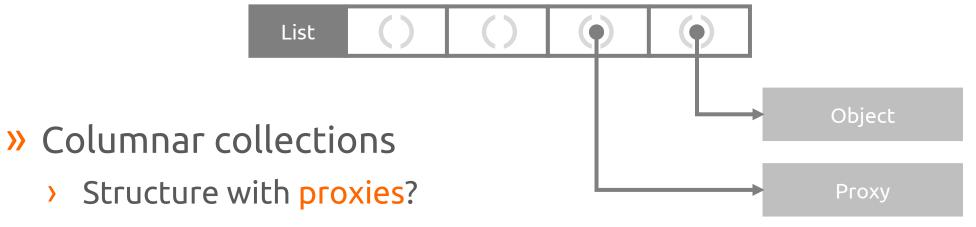
```
n = 0
col = LineItem.columns['quantity']
while n < size:

value = col[n]
total = total + value
n += 1</pre>
```

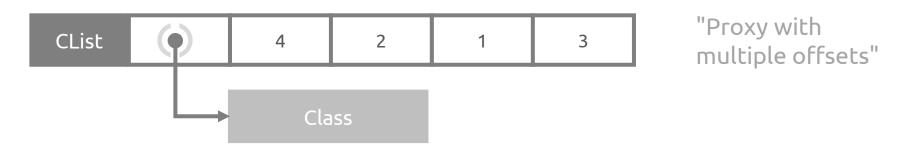
Changes to Collections



» Traditional collections: structure with pointers (addresses)

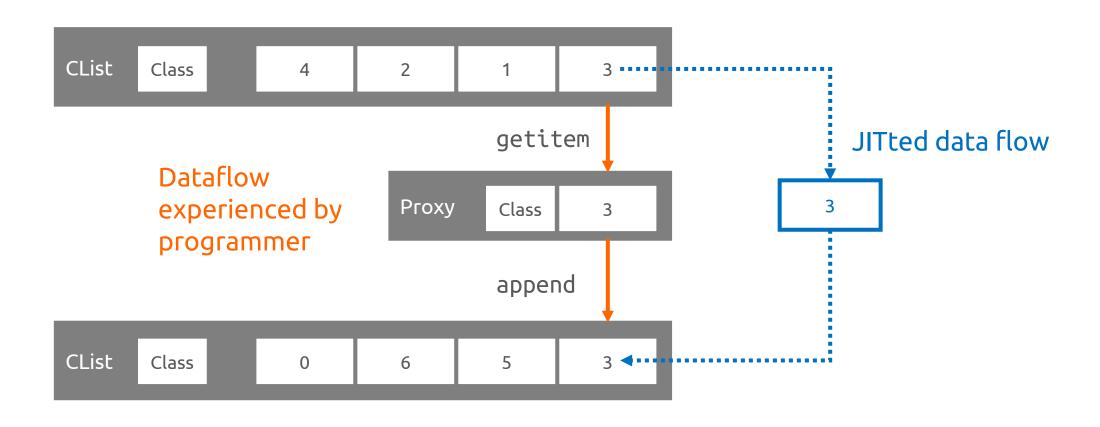


Structure with offsets if items of single class!



Collections and Allocation Removal

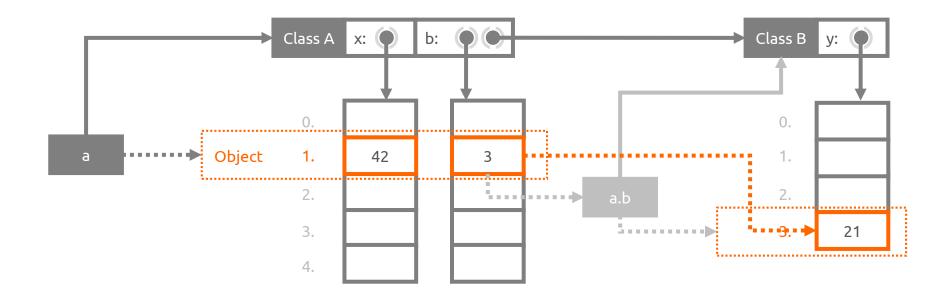




Associations



» Inspired by foreign keys



» Allocation Removal eliminates intermediate proxies:

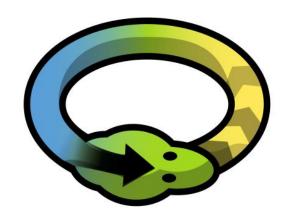
```
> a.b.y => y_column[b_column[offset]] ("Simple Join")
```

Our Prototype



PyPy

- Fast Python interpreter
- Tracing JIT-compiler
- Allocation Removal



Our Prototype "Obsidian"

- Implemented on top of PyPy (mostly library code)
- Optimized proxies and collections for PyPy's JIT

Evaluation Setup



Hypotheses

- Columnar objects outperform traditional objects on analytical code
- > Columnar objects are outperformed by commercial in-memory DBs

Competitors

- > Baseline: PyPy (as idiomatic Python code)
- > PyPy with columnar objects (as idiomatic Python code)
- Commercial in-memory database (as stored procedure)

Platform: 2x 6-core Intel Xeon E5 @ 2.3 GHz (24 threads) | 128 GB RAM | PyPy 2.5.0 | SLES 11.2

Evaluation Scenarios



ATP: Available-to-Promise Candidate Selection*

Determine a conflic-free delivery schedule for a set of orders given fixed incoming and outgoing stock changes (always stock ≥ 0)

KM: Kaplan-Meier Estimator*

Approximate a survival function over time given censored records

Elo: Chess-Player Ranking

Rank players given only their match outcomes

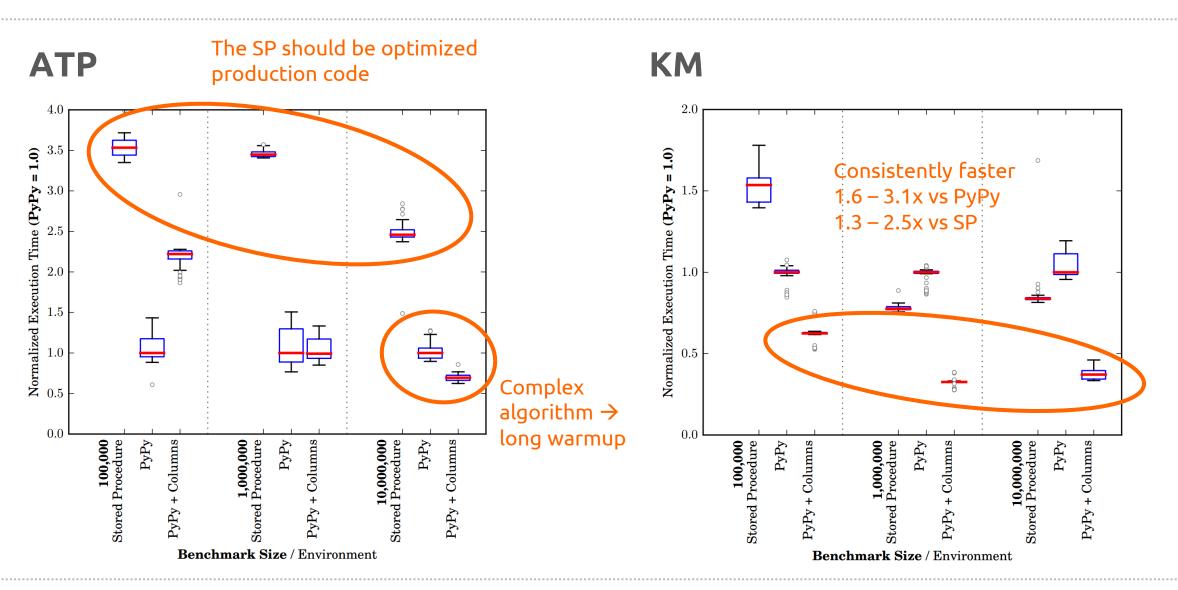
Balance:

> Compute how long an account was overdrawn given a set of transactions

*) professionally optimized stored procedures were provided by the database vendor

Results

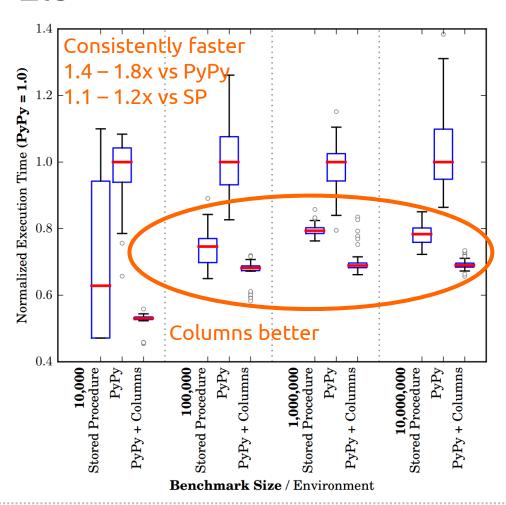




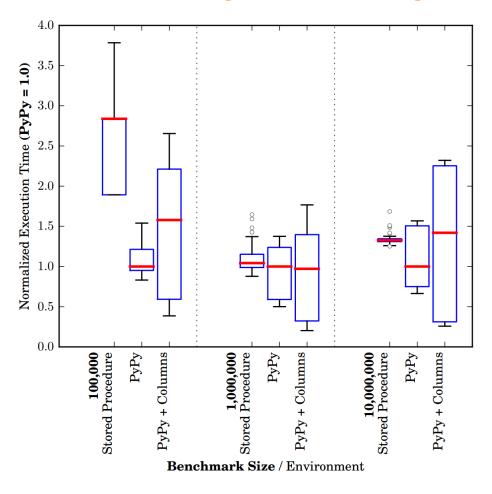
Results



Elo



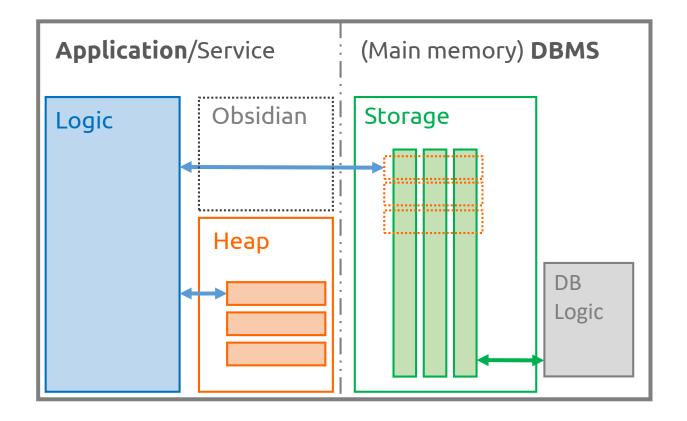
Balance No significance due to high variance



Future Work

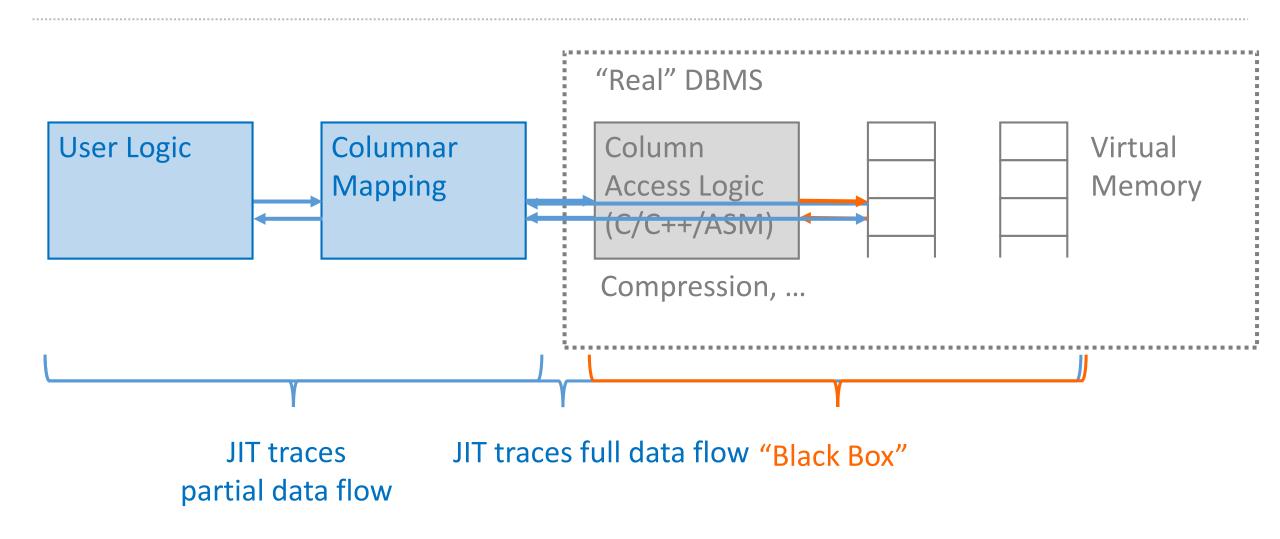


- » Run obsidian on the same columns as an in-memory database
 - Ongoing research in our group with promising preliminary results



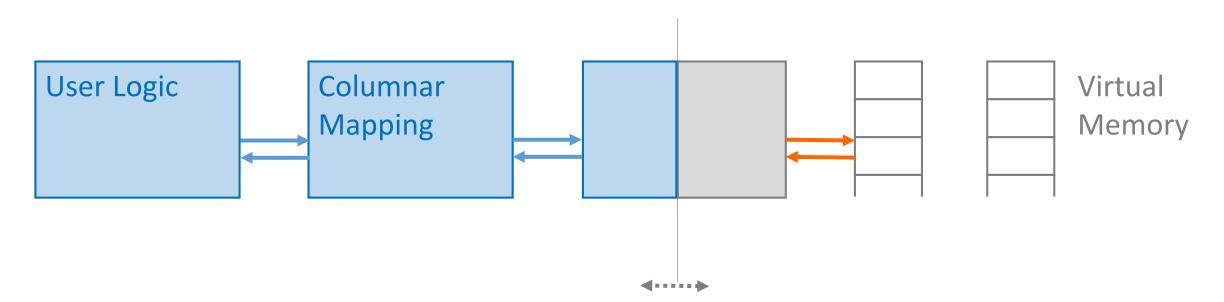
"Jit-compiling" to a Real Database





"Jit-compiling" to a Real Database





Where should we place the boundary?

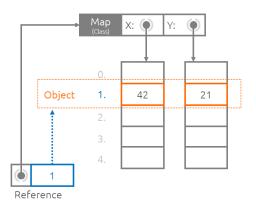
- ← Guide JIT to better "inline" C/C++ code
- → Re-implement DB in Python to allow JITting

Ongoing research by Johannes Henning, HPI

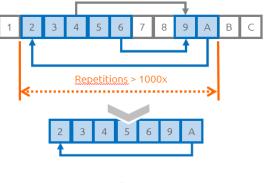
Conclusion



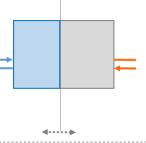
» Columns can bring performance benefits to dynamic language implementations in analytical scenarios



>> Tracing JIT compilers and columns syngergize well



» Unexplored opportunities in the database domain





Stored Procedures



```
CREATE FUNCTION "membership weight with skipping"
   [ ( "area id p" "area"."id"%TYPE, 23 Void removeCandidates(CANDIDATE T & candidates,
     "member id p" "member"."id"%TYPE, 24
                                                                      QTY T accumulated)
      "skip member ids p" INT4[] ) -- "m 25 -{
       RETURNS INT4
                                                    COL DATE T dates = candidates.getColumn<DATE T>("DDATE");
                                          26
       LANGUAGE 'plpgsgl' STABLE AS $$
                                                    COL QTY T values = candidates.getColumn<QTY T>("QTY01");
                                          27
       DECLARE
                                                    Size currSize = Size(dates.getSize());
                                                   QTY T lastValue;
        "sum v" INT4;
         "delegation row" "area delegatio 30
                                                   QTY T zero;
       BEGIN
                                                    while ((accumulated < zero) && (currSize > 0z)) {
         "sum v" := 1;
11
                                                        lastValue = values[currSize - 1z];
         FOR "delegation row" IN
12
                                                       if ((lastValue + accumulated) <= zero)</pre>
           SELECT "area delegation".*
           FROM "area delegation" LEFT JO 35
                                                            accumulated = accumulated + lastValue:
           ON "membership". "area id" = "a 36
                                                            if (currSize == 1z) { currSize = 0z; } else { currSize = currSize - 1z; }
           AND "membership"."member id" = 37
                                                            dates.setSize(currSize);
           WHERE "area delegation". "area 38
17
                                                            values.setSize(currSize);
           AND "area delegation". "trustee 39
           AND "membership". "member id" I 40
                                                        else
20
         LOOP
                                                            values[currSize - 1z] = values[currSize - 1z] + accumulated;
           IF NOT
             "skip member ids p" @> ARRAY 43
                                                            accumulated = zero;
           THEN
             "sum v" := "sum v" + "member 45
               "area id p",
               "delegation row"."truster 47
               "skip member ids p" || "de 48
```

Interacting with a Database



Object-relational mapping

- + Object-oriented abstractions
- Limited by underlying protocol (SQL, libpg, ...)
- Copying

Stored Procedures

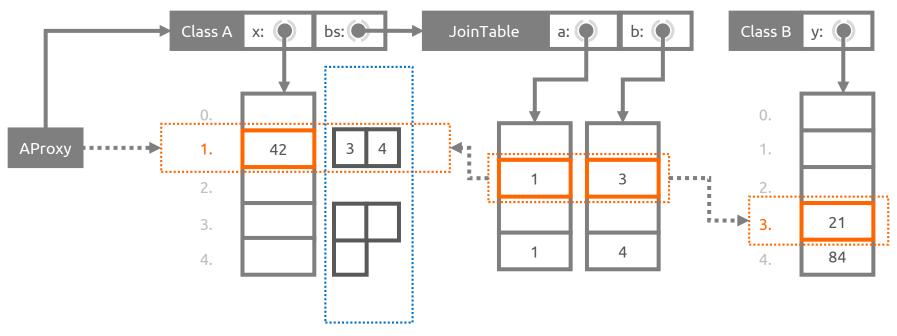
- + Performance
- Split logic (split tooling, split lifecycles)
- Technical abstactions (obscuring domain logic)

It may be a good idea to stay inside a single execution environment

Associations of higher Multiplicity



» Indexed Join-tables



Quasi-Index [Optional]

Configurable: Consists of per-instance Lists/Sets/Dicts

Collection Operations



» Many languages have special collection operators:

```
> Python: reduce, map, [f(a, b) for a in A for b in B]...
> ST-80: collection inject:into:, c collect:, c gather:, ...
> C#: collection.Reduce(func), c.Select(func), ...
```

» Example in Python:

```
sum(i.quantity * i.unit_price
    for i in order.items
    for order in customer.orders
    if order.year >= 2015)
```

Optimizing Collection Operations



» Defer evaluation

- Apply optimizations over the full operation
- > Prevent intermediate proxies
- > Leave "Column World" as late as possible

» Infer result types

> Allows to allocate result columns instead of ordinary objects

Deferred Evaluation: Plan Construction



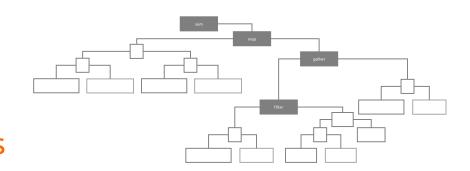
```
sum(i.quantity * i.unit_price
                                                           Python OpCodes
    for i in order.items
    for order in customer.orders
    if order.year >= 2015)
                                                                       Abstract Interpretation
                                            sum
                                                        map
                                                                      gather
                                    arg[0]
            arg[0]
                       quantity
                                              unit price
                                                              filter
                                                                                    arg[0]
                                                                                  2015
                                                                  arg[0]
                                            customer
                                                        orders
                                                                            vear
```

Plan Optimization



Type inference

- LINQ type system by Erik Meijer (handles Python's dynamicity well)
- Create anonymous classes with result columns
- Warn user on failure, continue un-optimized



Tree transformations

- Move filters down the hierarchy
- Replace gather by relational (hash-)join
- **>**

Compile to new OpCode and run if needed

Results: Exact Timings



benchmark		platform timings [ms]			speedup	
	size	PyPy	SP	Col.	vs. PyPy	vs. SP
ATP	10 000	1.32	7.0	28.05	0.05 [0.04 - 0.05]	0.25 [0.25 - 0.29]
	100 000	21.79	77.0	48.41	0.45 [0.44 - 0.46]	1.59 [1.56 - 1.62]
	1000000	235.71	751.0	228.91	1.03 [0.93 - 1.16]	3.28 [3.2 - 3.38]
	10 000 000	2739.53	6736.5	1902.78	1.44 [1.4 - 1.49]	3.54 [3.48 – 3.61]
KM	10 000	0.59	4.0	2.99	0.2 [0.2 - 0.2]	1.34 [1.33 - 1.52]
	100 000	28.65	44.0	17.89	1.6 [1.59 - 1.62]	2.46 [2.4 - 2.52]
	1000000	535.12	415.0	174.31	3.07 [3.06 – 3.08]	2.38 [2.37 - 2.4]
	10 000 000	4393.76	3682.5	1631.58	2.69 [2.53 - 2.91]	2.26 [2.13 - 2.4]
Elo	10 000	6.36	4.0	3.38	1.88 [1.8 - 1.95]	1.18 [1.18 - 1.48]
	100 000	41.54	31.0	28.32	1.47 [1.43 - 1.51]	1.09 [1.09 - 1.12]
	1 000 000	359.06	285.0	247.49	1.45 [1.41 - 1.47]	1.15 [1.14 - 1.16]
	10 000 000	3506.49	2747.5	2418.84	1.45 [1.41 - 1.49]	1.14 [1.12 - 1.15]
Balance	10 000	0.11	0.0	0.16	0.7 [0.58 - 0.88]	O.O [0.0 - 0.0]
	100 000	1.06	3.0	1.67	0.63 [0.54 - 0.94]	1.8 [1.53 - 2.66]
	1 000 000	18.22	19.0	17.7	1.03 [0.61 - 1.91]	1.07 [0.81 - 2.03]
	10 000 000	119.85	159.0	170.17	0.7 [0.43 - 3.48]	0.93 [0.67 - 4.0]